## Object Constraint Language Specification

This chapter introduces and defines the Object Constraint Language (OCL), a formal language to express side-effect-free constraints.

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### 7.1 Overview

This chapter introduces and defines the Object Constraint Language (OCL), a formal language used to express constraints. These typically specify invariant conditions that must hold for the system being modeled. Note that when the OCL e xpressions are evaluated, they do not have side effects (i.e., their evaluation cannot alter the state of
the corresponding executing system). In addition to specifying invariants of the UML metamodel, UML modelers can use OCL to specify application-specific constraints in their models.

OCL is used in the UML Semantics chapter to specify the well-formedness rules of the metaclasses com prising the UML metamodel. A well-f ormedness rule in the static semantics chapters in the UML Semantics section normally co ntains an OCL expression, specifying an invariant for the associated metaclass. The grammar for OCL is specified at the end of this chapter. A parser generated from this gram mar has correctly p arsed all the con straints in the UML Semantics section, a p rocess which improved the correctness of the specifications for OCL and UML.

### 7.1.1 Why OCL?

A UML diagram, such as a class diag ram, is typically not refined enough to provide all the relevant aspects of a specification. There is, amo ng other things, a need to describe additional constraints ab out the ob jects in the model. Such constraints are ofte $n$ described in natural language. Practice has $s$ hown that this will a lways result in ambiguities. In order to write unambiguous con straints, so-calle d formal lang uages have been developed. The disadvantage of traditional formal languages is that they are usable to persons with a str ing mathematical back ground, but difficult for the a verage business or sy stem modeler to use.

OCL has been de veloped to fill this gap. It is a formal language that remains easy to read and write. It has been developed as a business modeling language within the IBM Insurance division, and has its roots in the Syntropy method.

OCL is a pure expression language; there fore, an OCL expression is $g$ uaranteed to be without side effect. When an OCL expression is evaluated, it sim ply returns a value. It cannot change anything in the model. This means that the state of the system will never change because of the evaluation of an OCL expression, even though an OCL expression can be used to specify a state change (e.g., in a post-condition).

OCL is not a prog ramming language; there fore, it is not possible to write program logic or flow control in OCL. You cannot invoke processes or activate non-query operations within OCL. Because OCL is a mo deling language in the first place, not everything in it is $p$ romised to be directly executable.

OCL is a ty ped language, so that each OCL expression has a ty pe. To be well formed, an OCL expression must conform to the type conformance rules of the language. For example, you cannot compare an Integer with a Str ing. Each Classifier defined within a UML model represents a distinct OCL ty pe. In add ition, OCL in cludes a set of supplementary predefined ty pes (these are described in Section 7.8, "Predefined OCL Types," on page 7-28).

As a specification language, all im plementation issues are out of scope and cannot be expressed in OCL.

The evaluation of an OCL expression is instantaneous. This means that the states of objects in a model cannot change during evaluation.

### 7.1.2 Where to Use OCL

OCL can be used for a n umber of different purposes:

- To specify invariants on classes and types in the class model
- To specify type invariant for Stereotypes
- To describe pre- and post con ditions on Op erations and Methods
- To describe Guards
- As a navigation language
- To specify constraints on operations

Within the UML Semantics chap ter, OCL is $u$ sed in the well- formedness rules as invariants on the metaclasses in the abstract syntax. In several places, it is also used to define 'additional' operations, which are used in the well-f ormedness rules.

### 7.2 Introduction

### 7.2.1 Legend

Text written in the courier typeface as shown below is an OCL expression.
'This is an OCL expression'
The context keyword introduces the context for the expression. The keywords inv, pre and post denote the stereotypes, respectively «invariant», «precondition», and «postcondition», of the constraint. The actual OCL expression comes after the colon.

```
context TypeName inv:
'this is an OCL expression with stereotype <<invariant>> in the
context of TypeName' = 'another string'
```

In the e xamples. the keywords of OCL are written in boldface in this document. The boldface has no formal meaning, but is used to make the expressions more readable in this document. OCL expressions are written using ASCII characters on ly.

Words in Italics within the main text of the paragraphs refer to parts of OCL expressions.

### 7.2.2 Example Class Diagram

The following diagram is used in the examples in this document.


Figure 7-1 Class Diagram Example

### 7.3 Connection with the UML Metamodel

### 7.3.1 Self

Each OCL expression is written in the context of an instance of a specific type. In an OCL expression, the reserved word self is used to refer to the contextual instance. For instance, if the context is Company, then self refers to an instance of Company.

### 7.3.2 Specifying the UML context

The context of an OCL expression within a UML model can be specified through a socalled context dec laration at the beginning of an OCL expression. The context declaration of the constraints in the follo wing section s is s hown.

If the con straint is shown in a diagram, with the proper stereotype and the dashed lines to connect it to its con textual element, there is no need for an e xplicit con text declaration in the test of the con straint. The context declaration is optional.

### 7.3.3 Invariants

The OCL expression can be part of an Invariant, which is a C onstraint stereotyped as an «invariant». When the invariant is associa ted with a Classifier, the la tter is referred to as a "type" in this chapter. An OCL e xpression is an invariant of the typ e and must be true for a ll instances of that type at an y time. (Note th at all OC L expressions that express invariants are of the type Boolean.)

For example, if in the context of the Company type in Figure 7-1, the following expression would specify an invariant that the nu mber of employees must always exceed 50:

```
self.numberOfEmployees > 50
```

where self is an instance of type Company. (We can view self as the object from where we start the expression.) This invariant holds for every instance of the Company type.

The type of the contextual in stance of an OC L expression, which is part of an invariant, is written with the context keyword, followed by the name of the type as follows. The label inv: declares the constraint to be an «invariant» constraint.

```
context Company inv:
    self.numberOfEmployees > 50
```

In most cases, the keyword self can be dropped because the context is clear, as in the above examples. As an alternative for self, a different name can be defined playing the part of self:

```
context c : Company inv:
    C.numberOfEmployees > 50
```

This invariant is equ ivalent to the p revious one.
Optionally, the n ame of the constraint m ay be written af ter the inv keyword, allowing the constraint to be referenced by name. In the following example the name of the constraint is enoughEmployees. In the UML metamodel, this name is an attribute of the metaclass C onstraint that is inh erited from ModelElement.

```
context c : Company inv enoughEmployees:
```

    c.numberOfEmployees > 50
    
### 7.3.4 Pre- and Postconditions

The OCL expression can be part of a Precondition or Postcondition, corresponding to «precondition» and «postcondition» stereotypes of Constraint associated with an Operation or Method. The contextual instan ce self then is an instance of the type that
owns the operation or method as af eature. The context declaration in OCL uses the context keyword, followed by the type and operation declaration. The stereotype of constraint is sh own by putting the labels 'pres:' and 'post:' before the actual Preconditions and Postconditions

```
context Typename::operationName(param1 : Type1, ... ): ReturnType
    pre : param1 > ...
post: result = ...
```

The name self can be used in the expression referring to the object on which the operation was called. The reserved word result denotes the result of the operation, if there is one. The n ames of the parameters (paraml) can also be used in the OCL expression. In the example diagram, we can write:

```
context Person::income(d : Date) : Integer
post: result = 5000
```

Optionally, the name of the precondition or po stcondition may be written after the are or post keyword, allowing the constraint to be referenced by name. In the following example the name of the precondition is parameter Ok and the name of the postcondition is resultOk. In the UML metamodel, these names are attributes of the metaclass Constraint that is inherited from ModelElement.

```
context Typename::operationName(param1 : Type1, ... ): ReturnType
    pre parameterOk: param1 > ...
    post resultok: result = ...
```


### 7.3.5 General Expressions

Any OCL expression can be used as the value for an attribute of the UML metac lass Expression or one of its subtypes. In that case, the semantics section describes the meaning of the expression.

### 7.4 Basic Values and Types

In OCL, a number of basic types are predefined and available to the modeler at all time. These predefined value types are independent of any object model and part of the definition of OCL.

The most basic value in OCL is a $v$ value of one of the $b$ asic ty pes. Some basic ty pes used in the examples in th is document, with co responding examples of their values, are shown in Table 7-1.

Table 7-1 Basic types

| type | values |
| :--- | :--- |
| Boolean | true, false |
| Integer | $1,-5,2,34,26524, \ldots$ |
| Real | $1.5,3.14, \ldots$ |
| String | 'To be or not to be...' |

OCL defines a number of operations on the predefined types. Table 7-2 gives some examples of the operations on the predefined types. See Section 7.8, "Predefined OCL Types," on page 7-28 for a com plate list of all operations.

Table 7-2 Operations on predefined types

| type | operations |
| :--- | :--- |
| Integer | $*,+,-, /$, abs |
| Real | $*,+,-, /$, floor |
| Boolean | and, or, xor, not, implies, if-then-else |
| String | toUpper, concat |

The complete list of operations provided for each type is described at the end of this chapter. Collection, Set, Bag and Sequence are basic ty pes as well. Their specifics will be described in the upcoming sections.

### 7.4. 1 Types from the UML Model

Each OCL expression is written in the context of a UML model, a number of classifiers (ty pes/classes, ...), the ir features and associations, and their generalizations. All classifiers from the UML model are types in the OCL expressions that are attached to the model.

### 7.4.2 Enumeration Types

As shown in the example diagram, new enumeration types can be defined in a model by using:

```
enum{ value1, value2, value3 }
```

The values of the e numeration can be used within expressions.
As there might be a name conflict with attribute nam es being equal to en u meration values, the usage of an enumeration value is expressed syntactically with an additional pound (\#) symbol prefixing the name of the value:

```
#value1
```

The type of an enumeration attribute is Enumeration, with restrictions on the values for the attribute.

### 7.4.3 Let Expression

Sometimes a sub-expression is used more than once in a constraint. The let expression allows one to define a variable that can be used in the constraint.

```
context Person inv:
    let income : Integer = self.job.salary->sum in
    if isUnemployed then
        income < 100
    else
        income >= 100
    endif
```


### 7.4.4 Type Conformance

OCL is a ty ped language and the basic value ty pes are organized in a type hierarchy. This her arch determines conformance of the different ty pes to each other. You cannot, for example, compare an Integer with a B oolean or a String.

An OCL expression in which all the types conform is a valid expression. An OC L expression in which the types don't conform is an invalid expression. It contains a type conformance error. A type type conforms to a type type 2 when an instance of type can be substituted at each place where an instance of type 2 is expected. The type conformance rules for types in the class diag grams are simple.

- Each type conforms to each of its supertypes.
- Type con formance is tran nsitive: if type conforms to type 2, and type 2 conforms to type, then type conforms to type.

The effect of this is that a type conforms to its supertype, and all the supertypes above. The type conformance rules for the value types are listed in Table 7-3.

Table 7-3 Type conformance rules

| Type | Conforms to/Is a subtype of |
| :--- | :--- |
| Set (T) | Collection(T) |
| Sequence (T) | Collection(T) |
| $\operatorname{Bag}(T)$ | Collection(T) |
| Integer | Real |

The conformance relation between the collection ty pes only holds if the y are collections of element types that conform to each other. See Section 7.5.14, "Collection Type Hierarchy and Type Conformance Rules," on page 7-20 for the complete conformance rules for collections.

Table 7-4 provides examples of valid and invalid expressions.
Table 7-4 Valid expressions

| OCL expression | valid | explanation |
| :--- | :--- | :--- |
| $1+2 * 34$ | yes |  |
| $1+$ 'motorcycle' | no | type Integer does not conform to type <br> String |
| $23 *$ false | no | type Integer does not conform to <br> Boolean |
| $12+13.5$ | yes |  |

### 7.4.5 Re-typing or Casting

In some circumstances, it is desirable to use a property of an object that is defined on a subtype of the current known type of the object. Because the property is not defined on the curr ant known ty pe, this results in a typ e conformance error.

When it is certain that the actual type of the object is the subtype, the object can be retyped using the op eration oclAsType(OclType). This operation results in the same object, but the known type is the argument OclType. When there is an object object of type Type 1 and Type 2 is another type, it is allowed to write:

```
object.oclAsType(Type2) --- evaluates to object with type Type2
```

An object can only be re-typed to one of its subtype; therefore, in the example, Type 2 must be a subtype of Type.

If the actual ty pe of the o bject is no ta sub type of the typ e to which it is re -typed, the expression is u defined (see Section 7.4.9, "Undefined Values," on page 7-10).

### 7.4.6 Precedence Rules

The precedence order for the operations, starting with highest precedence, in OCL is:

- @ ire
- dot and arrow operations: '.' and '->'
- unary 'not' and unary minus '-'
- '*' and '/'
- '+' and binary '-'
- 'if-then-else-endif'
- '<', '>', '<=', '>='
- '=’, '<>’
- 'and', 'or' and 'xor'
- 'implies'

Parentheses '(' and ')' can be used to ch ange precedence.

### 7.4.7 Use of Infix Operators

The use of infix operators is allowed in OCL. The operators '+', '-', '*’. ‘/', '<‘, ‘>’, '<>' '<=' '>=' are used as infix operators. If a type defines one of those operators with the correct signature, they will be $u$ sed as in fix operators. The expression:

$$
a+b
$$

is conceptually equal to the expression:

$$
a .+(b)
$$

that is, invoking the ' + ' operation on a with b as the parameter to the operation.
The infix operators defined for a type must have exactly one parameter. For the infix operators '<‘, '>’, ‘<=', '>=', ‘<>’, 'and, 'or', and 'xor' the return ty pe must be Boolean.

### 7.4.8 Comment

Comments in OCL are written n following two successive dashes (minus signs).
Everything immediately following the two dashes up to and including the end of line is part of the comment. For example:
-- this is a comment

### 7.4.9 Undefined Values

Whenever an OCL expression is being evaluated, there is a possibility that one or more of the queries in the expression are undefined. If this is the case, then the complete expression will be undefined.

There are two exceptions to this for the Boolean operators:

1. True OR-ed with anything is True
2. False AND-ed with anything is False

The above two rules are valid irrespective of the order of the arguments and the above rules are valid whether or no $t$ the $v$ blue of the other sub -expression is kn own.

### 7.5 Objects and Properties

OCL expressions can refer to C lassifiers (e.g., types, classes, interfaces, associations (acting as typ es) and datatypes). Also all attr ibutes, association-ends, methods, and operations witho ut side-effects that are defined on the se types can be used. In a class model, an operation or method is defined to be side-effect-free if the isQuery attribute of the operations is true. For the $p$ urpose of this $d$ ocument, we will $r$ efer to attributes, association-ends, and side-effect-free methods and operations as being properties. A property is one of:

- an Attribute
- an Ass ociationEnd
- an Operation with isQuery being true
- a Method with isQuery being true


### 7.5.1 Properties

The value of a property on an object that is de fined in a class dia gram is specified by a dot followed by the name of the pro perty.

```
context AType inv:
    self.property
```

If self is a reference to an object, then self.property is the value of the property property on self.

### 7.5.2 Properties: Attributes

For example, the age of a Person is written as self.age:

```
context Person inv:
    self.age > 0
```

The value of the subexpression self.age is the value of the age attribute on the particular in stance of Per son identif ied by self. The type of this subexpression is the type of the attr ibute age, which is the basic ty pe Inte ger.

Using attributes and operations defined on the basic value types, we can express calculations over the class model. For example, a business rule might be "the age of a Person is always greater than zero." This can be stated as shown in the invariant above.

### 7.5.3 Properties: Operations

Operations may have parameters. For example, as shown earlier, a Person object has an income expressed as a function of the date. This operation would be accessed as follows, for a Person aPerson and a date aDate:

```
aPerson.income(aDate)
```

The operation itself could be defined by a postcondition constraint. This is a constraint that is stereotyped as «postcondition». The object that is returned by the operation can be referred to by result. It takes the following form:

```
context Person::income (d: Date) : Integer
    post: result = age * 1000
```

The right-hand-side of this definition may refer to the operation being defined (i.e., the definition may be recursi ve) as long as the recursion is $n$ ot infinite. The ty pe of result is the return type of the operation, which is Integer in the above example.

To refer to an operation or a method that doesn't take a parameter, parentheses with an empty argument list ar e mandatory:

```
context Company inv:
    self.stockPrice() > 0
```


### 7.5.4 Properties: Association Ends and Navigation

Starting from a specific object, we can navigate an association on the class diagram to refer to other objects and their properties. To do so, we $n$ avigate the asso ciation by using the opposite as sociation-end:

```
object.rolename
```

The value of this expression is the set of objects on the other side of the rolename association. If the multiplicity of the ass ociation-end has a maximum of one (" $0 . .1$ " or " 1 "), then the value of this expression is an object. In the example class diagram, when we start in the con text of a Company (i.e., self is an instance of Company), we can write:

```
context Company
    inv: self.manager.isUnemployed = false
    inv: self.employee->notEmpty
```

In the first invariant self.manager is a Person, because the multiplicity of the association is one. In the secon d invariant self.employee will e valuate in a Set of Persons. By default, navigation will result in a Set. When the association on the Class Diagram is adorned with \{ordered\}, the navigation results in a Sequence.

Collections, like Sets, B ags, and Sequences are predefined types in OCL. They have a large number of predefined operations on them. A property of the collection itself is accessed by using an arrow '->' followed by the name of the property. The following example is in the context of a person :

```
context Person inv:
    self.employer->size < 3
```

This applies the size property on the Set self.employer, which results in the number of employers of the Per son self.

```
context Person inv:
```

```
self.employer->isEmpty
```

This applies the isEmpty property on the Set self.employer. This evaluates to true if the set of employers is empty and false otherwise.

### 7.5.4.1 Missing Rolenames

When a rolename is missing at one of the ends of an association, the name of the type at the assoc ration end, starting with a lowercase character, is used as the $r$ olename. If this results in an ambiguity, the rolename is mandatory. This is the case with unnamed rolenames in reflexive associations. If the rolename is ambiguous, then it can ot be used in OCL.

### 7.5.4.2 Navigation over Associations with Multiplicity Zero or One

Because the multiplicity of the role manager is one, self.manager is an object of type Person. Such a single object can be used as a Set as well. It then behaves as if it is a Set containing the single object. The usage as a set is done through the arrow followed by a property of Set. T his is sh own in the follow wing example:

```
context Company inv:
    self.manager->size = 1
```

The sub-expression self.manager is used as a Set, because the arrow is used to access the size property on Set. This expression evaluates to true.

```
context Company inv:
    self.manager->foo
```

The sub-expression self.manager is used as Set, because the arrow is used to access the foo property on the Set. This expression is incorrect, because foo is not a defined property of Set.

```
context Company inv:
    self.manager.age> 40
```

The sub-expression self.manager is used as a Person, because the dot is used to access the age property of Person.

In the case of an optional ( $0 . .1$ multiplicity) association, this is especially useful to check whether there is an object or not when navigating the association. In the example we can write:

```
context Person inv:
    self.wife->notEmpty implies self.wife.sex = #female
```


### 7.5.4.3 Combining Properties

Properties can be combined to make more complicated expressions. An important rule is that an OCL expression always evaluates to a specif ic object of a specific type. After obtaining a result, one can al ways apply another property to the result to get a ne w result value. Therefore, each OCL expression can be read and evaluated left-to-right.

Following are some invariants that use combined properties on the example class diagram:
[1] Married people are of age $>=18$

```
context Person inv:
    self.wife->notEmpty implies self.wife.age >= 18 and
    self.husband->notEmpty implies self.husband.age >= 18
```

[2] A company has at most 50 employees

```
context Company inv:
    self.employee->size <= 50
```


### 7.5.5 Navigation to Association Classes

To specify navigation to association classes (Job and Marriage in the example), OCL uses a dot and the name of the association class starting with a lowercase character:

```
context Person inv:
    self.job
```

The sub-expression self.job evaluates to a Set of all the jobs a person has with the companies that are his/her employer. In the ca se of an a ssociation class, th ere is no explicit rolename in the class d iagram. The name job used in this navigation is the name of the association class starting with a lowercase character, similar to the way described in the sectio $n$ "Missing Rolenames" above.

In case of a recursive association, that is an association of a class with itself, the name of the association class alo ne is not enough. We need to distinguish the direction in which the asso ciation is $n$ avigated as well as the name of the ass ociation class. Take the follo wing model as an example.


Figure 7-2 Navigating recursive association classes

When navigating to an association class such as employeeRanking there are two possibilities depen ding on the direction. For instance, in the above example, we $m$ ay navigate towards the employees end, or the bosses end. By using the name of the association class alon e, these two options cann ot be distin guished. To make the distinction, the rolename of the direction in wh ich we want to navigate is add ed to the association class name, enclosed in square brackets. In the expression

```
context Person inv:
    self.employeeRanking[bosses]->sum > 0
```

the self.employeeRanking[bosses] evaluates to the set of EmployeeRankings belonging to the collection of bosses. And in the expression

```
context Person inv:
    self.employeeRanking[employees]->sum > 0
```

the self.employeeRanking[employees] evaluates to the s et of EmployeeRankings belonging to the collectio $n$ of employees. The unqualified use of the association class name is not allowed in such a recursive situation. Thus, the following example is invalid:

```
context Person inv:
    self.employeeRanking->sum > 0 -- INVALID!
```

In a non-recursive situation, the associa tion class nam e alone is en ough, although the qualified version is allowed as well. Therefore, the examples at the start of this section could also be written as:

```
context Person inv:
    self.job[employer]
```


### 7.5.6 Navigation from Association Classes

We can navigate from the asso ciation class itself to th e objects th at participate in the association. This is d one using the dot-n otation and the role-names at the asso ciationends.

```
context Job
    inv: self.employer.numberOfEmployees >= 1
    inv: self.employee.age > 21
```

Navigation from an as sociation class to o ne of the objects on the association will always deliver exactly one object. This is a result of the definition of AssociationClass. Therefore, the result of this navigation is exactly one object, although it can be used as a Set using the arrow (->).

### 7.5.7 Navigation through Qualified Associations

Qualified asso ciations use o ne or m ore qualifier attributes to select the o bjects at the other end of the ass ociation. To navigate them, we can add the values for the qualifiers to the navigation. This is d one using square brackets, following the role- name. It is permissible to leave out the qualifier values, in which case the result will be all o bjects at the other end of the ass ociation.

```
context Bank inv:
    self.customer
```

This results in a Set(Person) containing all customers of the Bank.

```
context Bank inv:
    self.customer[8764423]
```

This results in one Person, having account number 8764423.
If there is more than one qualifier attribute, the values are separated by commas, in the order which is specified in the UML class model. It is $n$ ot permissible to partially specify the qualifier attribute values.

### 7.5.8 Using Pathnames for Packages

Within UML, different ty pes are organized in packages. OCL p rovides a way of explicitly referring to types in other packages by using a package-pathname prefix. The syntax is a pac kage name, followed by a do uble colon:

```
Packagename::Typename
```

This usage of pathnames is transitive and can also be used for packages within packages:

```
Packagename1::Packagename2::Typename
```


### 7.5.9 Accessing overridden properties of supertypes

Whenever properties are redefined within a type, the property of the supertypes can be accessed using the oclAsType() operation. Whenever we have a class B as a subtype of class A, and a property p 1 of both A and B, we can write:

```
context B inv:
    self.oclAsType(A).p1 -- accesses the p1 property defined in A
    self.p1 -- accesses the p1 property defined in B
```

Figure 7-3 shows an example where such a construct is needed.


Figure 7-3 Accessing Overridden Properties Example

In this model fragment there is an am biguity with the OCL expression on Dependency:

```
context Dependency inv:
    self.source <> self
```

This can either mean normal association navigation, which is inherited from ModelElement, or it $m$ ight also mean navigation through the dotted line as an association class. Bo th possible navigations use the same role-name, so th is is al ways ambiguous. Using oclAsType() we can distinguish between them with:

```
context Dependency
    inv: self.oclAsType(Dependency).source
    inv: self.oclAsType(ModelElement).source
```


### 7.5.10 Predefined properties on All Objects

There are several properties that apply to all objects, and are predefined in OCL. These are:

```
oclIsTypeOf(t : OclType) : Boolean
oclIsKindOf(t : OclType) : Boolean
oclInState(s : OclState) : Boolean
oclIsNew : Boolean
oclAsType(t : OclType) : instance of OclType
```

The operation is oclTypeOf results in true if the type of self and $t$ are the same. For example:

```
context Person
    inv: self.oclIsTypeOf( Person ) -- is true
    inv: self.oclIsTypeOf( Company) -- is false
```

The above property deals with th e direct type of an object. The oclIsKindOf property determines whether $t$ is either the direct type or one of the supertypes of an object.

The operation oclInState results in true if the object is in the state $s$. Values for $s$ are the names of the states in the statema chine(s) attach ed to the Classifier of object. For nested state $s$ the staten ames can be combined using the $::$.


Figure 7-4 Statemachine Example
In the previous example of statem achine, values for $s$ can be On, Off, Off::Standby, Off: : NoPower. If the classifier of object has the above associated statemachine, valid OCL expressions are:

```
object.oclInState(On)
object.oclInState(Off)
object.oclInstate(Off::Standby)
object.oclInState(Off:NoPower)
```

If there are multiple statemachines attached to the object's classifier, then the statename can be prefixed with the name of the statemachine containing the state and the double semicolon ::, as with nested states.

The operation ocllsNew evaluates to true if, used in a postcondition, the object is created during performing the op eration (i.e., it d idn't exist at precondition time).

### 7.5.11 Features on Classes Themselves

All properties discussed until now in OCL are properties on instances of classes. The types are either pre defined in OCL or def ined in the class $m$ odel. In OCL, it is also possible to use features defined on the ty pes/classes themselves. These are, for example, the class-scoped features defined in the class model. Furthermore, several features are predefined on each ty pe.

A predefined feature on each type is allInstances, which results in the Set of all instances of the type in existence at the specific time when the expression is evaluated. If we want to make sure that all instan ces of Person have unique names, we can write:

```
context Person inv:
    Person.allInstances->forAll(p1, p2 |
```

```
p1 <> p2 implies p1.name <> p2.name)
```

The Person.allInstances is the set of all persons and is of type Set(Person). It is the set of all persons that e xist at the s napshot in time that the expression is e valuated.

Note - The use of allInstances has some problems and its use is discouraged in most cases. The first problem is best explained by looking at the types like Integer, Real and String. For these types the meaning of allInstances is undefined. What does it mean for an Integer to exist? The evaluation of the expression Integer.allInstances results in an infinite set and is therefore undefined within OCL. The second problem with allInstances is that the existence of objects must be considered within some overall context, like a sy stem or a model. This overall con text must be defined, which is not done within OCL. A recommended style is to model the overall contextual system explicitly as an ob ject within the st stem and navigate from that object to its containing instances without using allInstances.

### 7.5.12 Collections

Single navigation results in a Set, combined navigation in a Bag, and navigation over associations adorned with $\{0$ ordered $\}$ results in a Sequence. Therefore, the collection types play an important role in OCL expressions.

The type Collection is predefined in OCL. The Collection type defines a large number of predefined operations to en able the OCL e xpression auth or (the modeler) to manipulate collections. Consistent with the definition of OCL as an expression language, collection operations never change collection s; is Query is always true. They may result in a collection, but rather than changing the original collection they project the result in to a new one.

Collection is an abstract type, with the concrete collection types as its sou btypes. OCL distinguishes three different collection types: Set, Sequ ence, and Bag. A Set is the mathematical set. It does not contain duplicate elements. A Bag is like a set, which may contain duplicates (i.e., the same element may be in a bag twice or more). A Sequence is like a Bag in which the elements are ordered. Both Bags and Sets have no order defined on them. Sets, Sequences, and Bags can be specified by a literal in OCL. Curly brackets sur round the elements of the collection, elements in the collection are written within, sep arated by commas. The typ e of the collection is written before the curly brackets:

```
Set { 1, 2, 5, 88}
Set { 'apple' , 'orange', 'strawberry' }
```

A Sequence:

```
Sequence { 1, 3, 45, 2, 3 }
Sequence { 'ape', 'nut' }
```

A bag:

```
Bag {1, 3 , 4, 3, 5 }
```

Because of the usefulness of a Seq uence of consecutive Integers, there is a sep arate literal to create them. The elements inside the curly brackets can be replaced by an interval specification, which consists of two expressions of type Integer, Int-exprl and Int-expr2, separated by '..'. This denotes all the I ntegers between the values of Intexpr1 and Int-expr2, including the values of Int-expr1 and Int-expr2 themselves:

```
Sequence{ 1..(6 + 4) }
Sequence{ 1..10 }
-- are both identical to
Sequence{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 }
```

The complete list of Co llection operations is described at the end of this chapter.
Collections can be specified by a literal, as described above. The only other way to get a collection is by navigation. To be more precise, the only way to get a Set, Sequence, or Bag is:

1. a literal, this will result in a Set, Sequence, or Bag:
```
Set {1, 2, 3, 5, 7, 11, 13, 17 }
Sequence {1, 2, 3, 5, 7, 11, 13, 17 }
Bag {1, 2, 3, 2, 1}
```

2. a navigation starting from a single object can result in a collection:
```
Company
    self.employee
```

3. operations on collection s may result in new collections:
```
collection1->union(collection2)
```


### 7.5.13 Collections of Collections

Within OCL, all Co llections of Collectio ns are flattened au tomatically; therefore, the following two expressions have the same value:

```
Set{ Set{1, 2}, Set{3, 4}, Set{5, 6} }
Set{ 1, 2, 3, 4, 5, 6 }
```


### 7.5.14 Collection Type Hierarchy and Type Conformance Rules

In addition to the ty pe conformance rules in S ection 7.4.4, "Type Conformance," on page 7-8, the following rules hold for a ll ty pes, in cluding the collection ty pes:

- The types Set (X), Bag (X) and Sequence (X) are all subtypes of Collection (X).

Type conformance rules are as follows for the collection types:

- Typel conforms to Type 2 when they are identical (standard rule for all types).
- Typel conforms to Type 2 when it is a subtype of Type 2 (standard rule for all types).
- Collection(Type1) conforms to Collection(Type2), when Type conforms to Type 2.
- Type conformance is transitive: if Type conforms to Type 2, and Type 2 conforms to Type, then Type conforms to Type (standard rule for all types).

For example, if Bicycle and Car are two separate subtypes of Transport:

```
Set(Bicycle) conforms to Set(Transport)
Set(Bicycle) conforms to Collection(Bicycle)
Set(Bicycle) conformsto Collection(Transport)
```

Note that $\operatorname{Set}$ (Bicycle) does not conform to Bag(Bicycle), nor the other way around. They are both subtypes of Collection(Bicycle) at the same le vel in the hierarchy.

### 7.5.15 Previous Values in Postconditions

As stated in Section 7.3.4, "Pere- and Postconditions," on page 7-5, OCL can be used to specify pere- and post-conditions on Operations and Methods in UML. In a postcondition, the expression can ref er to two sets of $v$ glues for each property of an object:

- the value of a property at the start of the operation or method
- the value of a property up on completion of the operation or method

The value of a prop erty in a po stcondition is the value up on completion of the operation. To refer to the value of a property at the start of the o peration, one has to postfix the property name with the keyword ' @ pere':

```
context Person::birthdayHappens()
    post: age = age@pre + 1
```

The property age refers to the prop erty of the instance of Person on which executes the operation. The property age@pre refers to the value of the property age of the Person that executes the op eration, at the start of the operation.

If the property has parameters, the '@ re' is postfixed to the propertyname, before the parameters.

```
context Company::hireEmployee(p : Person)
post: employees = employees@pre->including(p) and
        stockprice() = stockprice@pre() + 10
```

The above operation can also be specified by a po stcondition and a pere condition together:

```
context Company::hireEmployee(p : Person)
    pre : not employee->includes(p)
    post: employees->includes(p) and
        stockprice() = stockprice@pre() + 10
```

When the pre-value of a property evaluates to an object, all further properties that are accessed of this object are the $n$ ew values (upon completion of the operation) of this object. So:

```
a.b@pre.c -- takes the old value of property b of a, say x
    -- and then the new value of c of x.
a.b@pre.c@pre -- takes the old value of property b of a, say x
    -- and then the old value of c of x.
```

The '@ pre' postfix is allowed only in OCL expressions that are part of a Postcondition. Asking for a current property of an object that has been destroyed during execution of the o peration results in Undefined. Also, re ferring to the previous value of an object that has been created during execution of the operation results in Undefined.

### 7.6 Collection Operations

OCL defines many operations on the collection types. These operations are specifically meant to enable a flexible and powerful way of projecting new collections from existing ones. The different constructs are described in the follo wing section s .

### 7.6.1 Select and Reject Operations

Sometimes an expression using operations and navigations delivers a collection, while we are intere sted only in a special subset of the collection. OCL has special constructs to specify a selection from a spec ific collection. These are the select and reject operations. The select specifies a subset of a collection. A selec $t$ is an operation on a collection and is specified using the ar row-syntax:

```
collection->select( ... )
```

The parameter of select has a special syntax that enables one to specify which elements of the collection we want to select. There are three different forms, of which the simplest one is:
collection->select ( boolean-expression )
This results in a collection that contains all the elements from collection for which the boolean-expression evaluates to true. To find the result of this e xpression, for each element in collection the expression boolean-expression is evaluated. If this evaluates to true, the element is includ ed in the re sult collection, otherwise no $t$. As an example, the following OCL expression specifies that the collection of all the employees older than 50 years is not empty:

```
context Company inv:
    self.employee->select(age > 50)->notEmpty
```

The self.employee is of type $\operatorname{Set}(\mathrm{Person})$. The select takes each person from self.employee and evaluates age > 50 for this person. If this results in true, then the person is in the result Set.

As shown in the pre vious example, the conte xt for the e xpression in the select argument is the element of the collection on which the select is in voked. Thus the age property is ta ken in the context of a person.

In the above example, it is impossible to refer explicitly to the persons themselves; you can only refer to properties of them. To refer to the persons themselves, there is a more general sy ntax for the select e xpression:

```
collection->select( v | boolean-expression-with-v )
```

The variable $v$ is called the iterator. When the select is evaluated, $v$ iterates over the collection and the boolean-expression-with-v is evaluated for each $v$. The $v$ is a reference to the object from the collection and can be used to refer to the objects themselves from the collection. The two examples below are identical:

```
context Company inv:
    self.employee->select(age > 50)->notEmpty
context Company inv:
    self.employee->select(p | p.age > 50)->notEmpty
```

The result of the complete select is the collection of persons $p$ for which the p.age > 50 evaluates to True. This amo unts to a su bset of self.employee.

As a final extension to the select syntax, the e xpected ty pe of the $v$ ariable $v$ can $b e$ given. The select no $w$ is written as:

```
collection->select( v : Type | boolean-expression-with-v )
```

The meaning of this is th at the ob jects in collection must be of type Type. The next example is identical to the previous examples:

```
context Company inv:
    self.employee.select(p : Person | p.age > 50)->notEmpty
```

The compete select syntax now looks like one of:

```
collection->select( v : Type | boolean-expression-with-v )
collection->select( v | boolean-expression-with-v )
collection->select( boolean-expression )
```

The reject operation is identical to the select operation, but with reject we get the subset of all the elements of the collection for which the expression evaluates to F alse. The reject syntax is identical to the select syntax:

```
collection->reject( v : Type | boolean-expression-with-v )
collection->reject( v | boolean-expression-with-v )
collection->reject( boolean-expression )
```

As an example, specify that the collection of all the employees who are not married is empty:

```
context Company inv:
    self.employee->reject( isMarried ) ->isEmpty
```

The reject o peration is a vailable in OCL for convenience, because each reject can be restated as a select with the n egated expression. Therefore, the following two expressions are identical:

```
collection->reject( v : Type | boolean-expression-with-v )
collection->select( v : Type | not (boolean-expression-with-v) )
```


### 7.6.2 Collect Operation

As shown in the previous section, the select and reject operations always result in a sub-collection of the original collection. When we want to specify a collection which is derived from some other collection, but which contains different objects from the original collection (i.e., it is not a sub-collection), we can use a collect operation. The collect operation uses the same syntax as the select and reject and is written as one of:

```
collection->collect( v : Type | expression-with-v )
collection->collect( v | expression-with-v )
collection->collect( expression )
```

The value of the reject operation is the collection of the results of all the evaluations of expression-with-v.

An example: specif y the colle ction of birthDates for all employees in the context of a company. This can be written in the context of a Co mpany object as on e of:

```
self.employee->collect( birthDate )
self.employee->collect( person | person.birthDate )
self.employee->collect( person : Person | person.birthDate )
```

An important issue here is that the resulting collection is not a Set, but a Bag. When more than one employee has the same value for birthDate, this value will be an element of the resulting Bag more than once. The Bag resu lting from the collect operation always has the same size as the origin al collection.

It is possible to make a Set from the Bag, by using the asSet property on the Bag. The following expression results in the Set of different birthDates from all employees of a Company:

```
self.employee->collect( birthDate ) ->asSet
```


### 7.6.2.1 Shorthand for Collect

Because navigation through many objects is very common, there is a shorthand notation for the collect that makes the OCL expressions more readable. Instead of

```
self.employee->collect(birthdate)
```

we can also write:

```
self.employee.birthdate
```

In general, when we ap ply a property to a collection of Objects, then it will automatically be interpreted as a collect over the members of the collection with the specified property.

For any propertyname that is def ind as a pr operty on the objects in a collection, the following two expressions are eden tical:

```
collection.propertyname
collection->collect(propertyname)
```

and so are these if the property is parameterized:

```
collection.propertyname(par1, par2, ...)
collection->collect(propertyname(par1, par2, ...)
```


### 7.6.3 ForAll Operation

Many times a constraint is needed on all elements of a collection. The forAll operation in OCL allows specifying a Boolean expression, which must hold for all objects in a collection:

```
collection->forAll( v : Type | boolean-expression-with-v )
collection->forAll( v | boolean-expression-with-v )
collection->forAll( boolean-expression )
```

This forAll expression results in a Boolean. The result is true if the boolean-expression-with-v is true for a ll elements of collection. If the boolean-expression-with$v$ is false for one or more $v$ in collection, then the com plate expression evaluates to false. For example, in the context of a company:

```
context Company
    inv: self.employee->forAll( forename = 'Jack' )
    inv: self.employee->forAll( p | p.forename = 'Jack' )
    inv: self.employee->forAll( p : Person | p.forename = 'Jack' )
```

These invariants evaluate to true if the forename feature of each employee is equal to 'Jack.'

The forAll operation has an extended variant in which more then one iterator is used. Both iterators will iterate over the com plate collection. Effectively this is a forAll on the Cartesian p roduct of the collection with itself.

```
context Company inv:
    self.employee->forAll( e1, e2 |
        e1 <> e2 implies e1.forename <> e2.forename)
context Company inv:
    self.employee->forAll( e1, e2 : Person |
        e1 <> e2 implies e1.forename <> e2.forename)
```

This expression evaluates to true if the forenam es of all employees are different. It is semantically equivalent to:

```
context Company inv:
    self.employee->forAll(e1 | self.employee->forAll (e2 |
    e1 <> e2 implies e1.forename <> e2.forename)))
```


### 7.6.4 Exists Operation

Many times one needs to know whether there is at least one element in a collection for which a con straint holds. The exists operation in OCL allows you to specify a Boolean expression which must hold for at lea st one object in a collection:

```
collection->exists( v : Type | boolean-expression-with-v )
collection->exists( v | boolean-expression-with-v )
collection->exists( boolean-expression )
```

This exists operation results in a Boolean. The result is true if the boolean-expression-with-v is true for at least one element of collection. If the boolean-expression-with-v is false for all $v$ in collection, then the complete expression evaluates to false. For example, in the context of a com pany:

```
context Company inv:
    self.employee->exists( forename = 'Jack' )
context Company inv:
    self.employee->exists( p | p.forename = 'Jack' )
context Company inv:
    self.employee->exists( p : Person | p.forename = 'Jack' )
```

These expressions evaluate to true if the forename feature of at least one employee is equal to 'Jack.'

### 7.6.5 Iterate Operation

The iterate operation is slightly more complicated, but is very generic. The operations reject, select, forAll, exists, collect, can all be described in terms of iterate.

An accumulation builds one value by iterating over a collection.

```
collection->iterate( elem : Type; acc : Type = <expression> |
    expression-with-elem-and-acc )
```

The variable elem is the iterator, as in the definitions of select and forAll. The variable acc is the accumulator. The accumulator gets an initial value <expression>.

When the iterate is evaluated, elem iterates over the collection and the expression-with-elem-and-acc is evaluated for each elem. After each evaluation of expression-with-elem-and-acc, its value is ass igned to acc. In this way, the value of acc is built up during the iteration of the collection. The collect op eration described in terms of iterate will look like:

```
collection->collect(x : T | x.property)
-- is identical to:
    collection->iterate(x : T; acc : T2 = Bag{} |
    acc->including(x.property))
```

Or written in Java-like pseu docode the result of the iter ate can be calculated as:

```
iterate(elem : T; acc : T2 = value)
{
    acc = value;
    for(Enumeration e = collection.elements() ;
e.hasMoreElements(); ){
            elem = e.nextElement();
            acc = <expression-with-elem-and-acc>
    }
}
```

Although the Java pseudo code uses a ' next element,' the iterate operation is defined for each collection type and the order of the iteration through the elements in the collection is not defined for Set and Bag. For a Sequ ence the order is the o rder of the elements in the sequence.

### 7.7 The Standard OCL Package

Each UML model that uses OCL constraints contains a predefined standard package called "U ML_OCL." This package is used by default in all oth er packages in the model to e valuate OCL expressions. This package contains all predefined OCL types and their features.

To extend the predefined OCL types, a modeler should define a separate package. The standard OCL package can be imported, and each OCL type can be extended with new features.

To specify that a package used the predefined OCL types from a user defined package instead of the standard package, the using package must define a Dependency with stereotype <<OCL_Types>> to the package which defines the extended OCL types.

A constraint on the user defined OCL package is that as a minimum all predefined OCL types with all of their features must be defined. The user defined package must be a proper extension to the standard OCL package.

### 7.8 Predefined OCL Types

This section contains all standard types defined within OCL, including all the properties defined on those types. Its signature and a description of its semantics define each property. Within the description $n$, the reser ved word 'result' is $u$ wed to refer to the value th at results from evaluating the property. In se veral places, po st con ditions are used to describe properties of the result. When there is more than one postcondition, all postconditions must be true.

### 7.8.1 Basic Types

The basic types used are Integer, Real, String, and Boolean. They are supplemented with OclExpression, OclType, and OclAny.

### 7.8.1.1 OclType

All types defined in a UML model, or pre-defined within OCL, have a type. This type is an instance of the OCL typ e called Ocl Type. Access to this ty pe allows the modeler limited access to the meta-level of the model. This can be useful for advanced modelers.

Properties of OclType, where the instance of OclType is called type.
type. name : String
The name of type.

## type.attributes: S et(String)

The set of names of the attributes of type, as they are defined in the model.

## type.associationEnds : Set(Str ing)

The set of names of the navigable associationEnds of type, as they are defined in the model.
type.operations : $\operatorname{Set}(S t r i n g)$
The set of names of the operations of type, as they are defined in the model.

## type.supertypes : S et(OclType)

The set of all direct supertypes of type.
post: type.allSupertypes->includesAll(result)
type.allSupertypes: Set(Oc IType)
The transitive closure of the set of all sup ertypes of type.

## type.allInstances : Set(typ e)

The set of all in stances of type and all its su btypes in existence at the snapshot at the time that the expression is evaluated.

### 7.8.1.2 OclAny

Within the OCL context, the type OclAny is the supertype of all types in the model and the basic predefined OCL type. The predefined OCL Collection types are not subtypes of OclAny. Properties of OclAny are available on each object in all OCL expressions.

All classes in a UML m odel in herit all p roperties defined on OclAny. To avoid name conflicts between properties in the model and the properties inherited from OclAny, all names on the properties of OclAny start with 'ocl.' Although theoretically there may still be name conflicts, the y can be avoided. One can also use the oclAsT ype() operation to explicitly refer to the OclAn y properties.

Properties of OclAn y, where the instan ce of OclAny is called object.

$$
\text { object }=(\text { object } 2: \text { OclAn } y): \text { B oolean }
$$

True if object is the same object as object2.

## object <> ( object2 : OclAn y) : B oolean

True if object is a different object from object2.
post: result $=$ no $t($ object $=$ ob ject2 $)$

```
object.oclIsKindOf(type : OclType) : B oolean
```

True if type is one of the types of object, or one of the supertypes (transitive) of the types of object.

## object.oclIsTypeOf(type : Oc lType) : B oolean

True if type is equal to one of the types of object.

```
object.oclAsType(type : Oc lType) : typ e
    Results in object, but of known type type.
    Results in Undefined if the actual type of object is not type or one of its
    subtypes.
    pre:object.oclIsKindOf(type)
    post: result = ob ject
    post: result.oclIsKindOf(type)
```

    object.oclInState(state: OclState) : B oolean
    Results in true if object is in the state state, otherwise results in f alse. The argument is a name of a state in the state machine corresponding with the class of object.

## object.oclIsNew : B oolean

Can only be used in a postcondition.
Evaluates to true if the object is created during performing the operation. I.e. it didn't exist at precondition time.

### 7.8.1.3 OclState

The type OclState is $u$ sed as a p arameter for the operation oclInState. There are no properties defined on OclState. One can only specify an OclState by using the name of the state, as it a ppears in a statem achine. These names can be fully qualified by the nested state $s$ and statemachine that c ontain them.

### 7.8.1.4 OclExpression

Each OCL expression itself is an object in the context of OCL. The type of the expression is OclExp ression. This ty pe and its p roperties are used to d efine the semantics of properties that take an expression as one of their parameters. For example; select, collect or forAll.

An OclExpression includes the optional iterator variable and type and the optional accumulator variable and type.

Properties of OclExpression, where the instance of OclExpression is called expression.
expression.evaluationType : OclType
The type of the ob ject that results from evaluating expression.

### 7.8.1.5 Real

The OCL type Real represents the mathematical concept of real. Note that Integer is a subclass of Real, so for each parameter of type Real, you can use an integer as the actual parameter.

Properties of Real, where the instance of Real is called $r$.

$$
\mathrm{r}=(\mathrm{r} 2: \mathrm{Re} \text { eal }): \text { B oolean }
$$

True if $r$ is equal to $r 2$.
r <> (r2 : Real) : Boolean
True if $r$ is not equal to $r 2$.
post: result $=\operatorname{not}(r=r 2)$
$r+(r 2: R e a l): R e a l$
The value of the ad dition of $r$ and $r 2$.
r-(r2: Real) : Real
The value of the subtraction of $r 2$ from $r$.

$$
\mathrm{r} *(\mathrm{r} 2: \text { Real }): \text { Real }
$$

The value of the multiplication of $r$ and $r 2$.

## r / (r2 : Real) : Real

The value of $r$ divided by $r 2$.
r.abs: Real

The absolute value of $r$. post: if $r<0$ th en result $=-r$ else $r$ exult $=r$ enif
r.floor: Integer

The largest integer which is less than or equal to $r$. post: (result $<=r$ ) and (result $+1>\mathrm{r}$ )

## r. round : Integer

The integer which is closest to $r$. When there are two such integers, the largest one.
post: $((\mathrm{r}-\mathrm{result})<\mathrm{r}) . \mathrm{abs}<0.5)$ or $((\mathrm{r}-\mathrm{result}) . \mathrm{abs}=0.5$ and $(\mathrm{result}>\mathrm{r}))$
r.max(r2 : Real) : Real

The maximum of $r$ and $r 2$.
post: if $r>=r 2$ then result $=r$ else $r$ exult $=r 2$ endif
r.min(r2: Real) : Real

The minimum of $r$ and $r 2$.
post: if $r<=r 2$ then result $=r$ else $r$ exult $=r 2$ endif
$r<(r 2:$ Real $)$ : Boolean
True if $r 1$ is less than $r 2$.
$r>(r 2$ : Real) : Boolean
True if $r 1$ is greater than $r 2$. post: result $=\operatorname{not}(\mathrm{r}<=\mathrm{r} 2)$
r <= (r2 : Real) : Boolean
True if r 1 is less than or equal to r 2 . post: result $=(r=r 2)$ or $(r<r 2)$
$r>=(r 2: R e a l): B$ oolean
True if $r 1$ is greater $r$ than or equal to $r 2$. post: result $=(r=r 2)$ or $(r>r 2)$

### 7.8.1.6 Integer

The OCL type Integer represents the $m$ athematical con cept of inte ger.
Properties of I integer, where the ins stance of Integer is called $i$.
$\mathrm{i}=(\mathrm{i} 2$ : Integer) : Boolean
True if $i$ is equal to $i 2$.
i + (in: Integer): Integer
The value of the addition of $i$ and $i 2$.
i - (in : Integer) : Integer
The value of the subtraction of $i 2$ from $i$.
i * (in : Integer) : Integer
The value of the multiplication of $i$ and $i 2$.
i / (in : Integer) : Real
The value of $i$ divided by $i 2$.

## i. abs: Integer

The absolute value of $i$. post: if $\mathrm{i}<0$ then resu lt $=-\mathrm{i}$ else r exult $=\mathrm{i}$ endif
i. div( in : Integer) : Integer

The number of times that $i 2$ fits completely within $i$. pres: in <> 0
post: if i $/ \mathrm{i} 2>=0$ then result $=(\mathrm{i} / \mathrm{i} 2)$.floor else result $=-((-\mathrm{i} / \mathrm{i} 2)$. floor $)$ endif
i.mod( in : Integer) : Integer

The result is $i$ modulo $i 2$.
post: result $=\mathrm{i}-(\mathrm{i} . \operatorname{div}(\mathrm{i} 2) * \mathrm{i} 2)$
i.max(i2 : Ante ger) : Integer

The maximum of $i$ an $i 2$. post: if $\mathrm{i}>=\mathrm{i} 2$ th en result $=\mathrm{i}$ else r exult $=\mathrm{i} 2$ end if
i.min(i2 : Ante ger) : Ante ger

The minimum of $i$ an $i 2$.
post: if $\mathrm{i}<=\mathrm{i} 2$ th en result $=\mathrm{i}$ else r esult $=\mathrm{i} 2$ end if

### 7.8.1.7 String

The OCL type String represents ASCII strings.
Properties of String, where the instance of String is called string.
string $=($ str ing $2:$ Str ing $):$ Bo olean
True if string and string 2 contain the same characters, in the sam e order.
string.size : I nteger
The number of characters in string.
string.concat(string2 : Str ing) : Str ing
The concatenation of string and string 2 . post: result.size $=$ strin g.size +s tring2.size
post: result.substring ( 1 , string. size ) $=$ string
post: result.substring(string.size +1 , result.size) $=$ string 2

## string.toUpper : String

The value of string with all lowercase characters converted to uppercase characters. post: result.size $=$ string.size
string.toLower: String
The value of string with all u ppercase characters converted to lowercase characters.
post: result.size $=$ strin g.size
string.substring(lower : Integer, upper : Integer) : String
The sub-string of string starting at character number lower, up to and including character number upper.

### 7.8.1.8 Boolean

The OCL type Boolean represents the common true/false values.
Features of Boolean, the instance of Boolean is called $b$.

$$
\mathrm{b}=(\mathrm{b} 2: \text { Boolean }): \text { Boolean }
$$

Equal if $b$ is the same as $b 2$.
b or (b2: Boolean) : Boolean
True if either $b$ or $b 2$ is true.
b xor (b2: Bo olean) : Bo olean
True if either $b$ or $b 2$ is true, but not both. post: ( b or b 2 ) and $\operatorname{not}(\mathrm{b}=\mathrm{b} 2)$
b and (b2: Boolean) : Boolean
True if both $b 1$ and $b 2$ are true.
not b: Boolean
True if $b$ is false.
post: if $b$ then result $=$ false else $r$ result $=$ true endif
b implies (b2: Boo lean) : Boo lean
True if $b$ is false, or if $b$ is true and $b 2$ is true. post: (not b) or (b and b2)
if $b$ then (expression $1:$ OclExp ression)
else (expressio n2: OclExpression) endif : expression1.evaluationType
If $b$ is true, the result is the value of evaluating expression $;$; otherwise, result is the value of evaluating expression 2.

### 7.8.1.9 Enumeration

The OCL type Enumeration represents the enumerations defined in an UML model.

Features of Enumeration, the instan ce of Enumeration is called enumeration.

```
enumeration = (enumeration2:B oolean) : B oolean
```

Equal if enumeration is the same as enumeration 2.

```
enumeration <> (enumeration2 : B oolean) : B oolean
```

Equal if enumeration is not the same as enumeration 2. post: result $=$ not $($ enumeration $=$ e numeration 2 $)$

### 7.8.2 Collection-Related Types

The following section s define the properties on collection s (i.e., these properties are available on Set, Bag, and Sequence). As defined in this section, each collection type is actually a term plate with one parameter. ' $T$ ' denotes the p arameter. A re al collection type is cr eated by substituting a ty pe for the T. So Set (Integer) and Bag (Person) are collection types.

### 7.8.2.1 Collection

Collection is the abstract supertype of all collection types in OCL. Each occurrence of an object in a collection is call ed an element. If an object occurs twice in a collection, there are two elements. This section defines the properties on Collections that have identical sem antics for all c ollection subtypes. Some properties may be def ind with the subtype as well, which means that there is an additional postcondition or a m ore specialized re turn value.

The definition of several common properties is different for ea ch subtype. These properties are not mentioned in this section.

Properties of Collection, where the instance of Collection is called collection.

## collection->size : I integer

The number of elem ments in the co lection collection. post: result $=$ collection- >iterate $($ elem; acc $:$ Integer $=0 \mid \mathrm{acc}+1)$

## collection->includes(object : OclAn y) : Bo olean

True if object is an element of collection, false otherwise. post: result $=($ collection $->\operatorname{count}($ object $)>0)$

```
collection->excludes(object:OclAny) : B oolean
```

True if object is not an element of collection, false oth erwise. post: result $=($ collection $->\operatorname{count}($ object $)=0)$
collection->count(object: OclAn y) : I nteger
The number of times that object occurs in the collection collection. post: result $=$ collection->iterate (elem; acc $:$ I nteger $=0 \mid$
if elem $=$ object then acc +1 else acc endif)
collection->includesAll(c2 : Co llection(T)) : Boo lean
Does collection contain all the elements of $c 2$ ?
post: result $=c 2->$ forAll $($ elem $\mid$ collection $->$ includes $($ elem $)$ )
collection->excludesAll(c2 : C ollection(T)) : Bo olean
Does collection contain none of the elements of $c 2$ ?
post: result $=\mathrm{c} 2->$ forAll $($ elem $\mid$ collection $->\operatorname{excludes}($ elem $))$
collection->isEmpty : B oolean
Is collection the empty collection? post: result $=($ collection $->$ size $=0)$
collection->notEmpty : Bo olean
Is collection not the empty collection?
post: result $=($ collection $->$ size $\langle>0)$

## collection->sum : T

The addition of all elem ments in collection. Elements must be of a type supporting the + operation. The + operation must take one parameter of type $T$ and be both associative: $(a+b)+c=a+(b+c)$, and commutative: $a+b=b+a$. Integer and Real fulfill this condition.

$$
\begin{aligned}
\text { post: result }= & \text { sole ction->iterate }(\text { elem } ; \text { acc }: T=0 \mid \\
& \text { acc }+ \text { elem })
\end{aligned}
$$

## collection->exists(expr : OclExp cession) : B oolean

Results in true if expr evaluates to true for at least one element in collection.

```
post: result = collection->iterate(elem; acc : B oolean = false 
```

    acc rexpr)
    collection->forAll(expr : Oc lExpression) : Boolean
Results in true if expr evaluates to true for e och element in collection; otherwise, result is false.

```
post: result = collection->iterate(elem; acc : B oolean = true |
    acc and expr)
```

collection->isUnique(expr : Oc IExpression) : B oolean
Results in true if expr evaluates to a different value for each element in collection; otherwise, result is false.
post: result $=$ collection->collect(expr)->forAll(e1, ez $\mid \mathrm{e} 1<>\mathrm{e} 2)$
collection->sortedB y(expr : Oc lExpression) : B oolean
Results in the Sequence containing all elements of collection. The element for which expr has the lowest value comes first, and so on. The type of the expr expression must have the < operation defined. The < operation must be transitive ie. if $\mathrm{a}<\mathrm{b}$ an $\mathrm{d} \mathrm{b}<\mathrm{c}$ th en $\mathrm{a}<\mathrm{c}$.
post:
collection->iterate(expr : OclExp ression) : expr.evaluationType
Iterates over the collection. See "Iterate Operation" on page 7-26 for a complete description. This is the $b$ asic collection operation with which the other collection operations can be described.

### 7.8.2.2 Set

The Set is the mathematical set. It contains elements without duplicates. Features of Set, the instance of Set is called set.
set->union(set2 : Set(T)) : Set(T)
The union of set and set 2 .
post: result->forAll(elem | set->includes(elem) or set2->inclu des(elem))
post: set->forAll(elem | result->includes(elem))
post: set2->forAll(elem | result->includ es(elem))
set->union(bag : Bag (T)) : Bag(T)
The union of set and bag.
post: result->forAll(elem |
result->count(elem) $=$ set->count(elem) + bag->count(elem))
post: set->forAll(elem | result->includes(elem))
post: bag->forAll(elem | result->includes(elem))
set $=(\operatorname{set} 2: \operatorname{Set}(T)):$ Bo olean
Evaluates to true if set and set 2 contain the same elements.
post: result $=($ set->forAll(elem $\mid$ set2->includes(elem) $)$ and
set2->forAll(elem | set->includes(elem)) )
set->intersection (set2 : Set(T)) : Set(T)
The intersection of set and set2 (i.e, the set of all elements that are in both set and set2).
post: result->forAll(elem | set-> includes(elem) and set2->includ es(elem))
post: set->forAll(elem | set2->inclu des(elem) $=$ result->includ es(elem))
post: set2->f orAll(elem | set->includes(elem) $=$ result->includ es(elem))
set->intersection (bag : Bag (T)) : Set(T)
The intersection of set and bag. post: result $=$ set->intersection $($ bag->asSet $)$
set $-(\operatorname{set} 2: \operatorname{Set}(T)): \operatorname{Set}(T)$
The elements of set, which are not in set 2 .
post: result->forAll(elem | set-> includes(elem) and set2->excludes(elem))
post: set->forAll(elem | result->includes(elem) = set2->e xcludes(elem))
set->including(object : T) : S et(T)
The set containing all elements of set plus object.
post: result->forAll(elem | set-> includes(elem) or (e lem = ob ject))
post: set->forAll(elem | result->includes(elem))
post: result->includes(object)
set->excluding(object : T ) : Set( T)
The set containing all elements of set without object.
post: result->forAll(elem | set-> includes(elem) and (elem <> object))
post: set->forAll(elem |result->includes(elem) $=($ object <> e lem $))$
post: result->excludes(object)
set->symmetricDifference(set2 : Set(T )) : Set(T)
The sets containing all the elements that are in set or set 2 , but not in both.
post: result->forAll(elem | set->includes(elem) x or set2->includes(elem))
post: set->forAll(elem | result->includes(elem) = set2 ->excludes(elem))
post: set2->forAll(elem | result->includ es(elem) = set->e xcludes(elem))
set->select(expr : OclE expression) : Set (T)
The subset of set for which expr is true.
post: result $=$ set->iterate $($ elem; acc $: \operatorname{Set}(T)=\operatorname{Set}\{ \} \mid$
if expr then acc->including(elem) else acc endif)
set->reject(expr : OclExp ression) : Set (T)
The subset of set for which expr is false. post: result $=$ set $->$ select $($ not expr)
set->collect(expr : OclExp ression) : Bag (expr.evaluationType)
The Bag of elements which results from applying expr to every member of set.
post: result $=$ set->iterate $($ elem; acc $: \operatorname{Bag}($ expr.evaluationType $)=\operatorname{Bag}\{ \} \mid$ acc->including(expr) )
set->count(object : T) : Integer
The number of occurrences of object in set.
post: result <= 1
set->asSequence : Seq uence(T)
A Sequence that contains all the elements from set, in undefined order.
post: result->forAll(elem | set-> includes(elem))
post: set->forAll(elem | result->count(elem) =1)
set->asBag : Bag (T)
The Bag that contains all the ell mints from set.
post: result->forAll(elem | set-> includes(elem))
post: set->forAll(elem | result->count(elem) $=1$ )

### 7.8.2.3 Bag

A bag is a collection with duplicates allowed. That is, one object can be an element of a bag many times. $T$ here is no ordering defined on the elements in abag.

Properties of Bag, where the in stance of Bag is called bag.
$\mathrm{bag}=(\mathrm{bag} 2: \mathrm{Bag}(\mathrm{T})):$ B oolean
True if bag and bag 2 contain the same elements, the same number of times. post: result $=($ bag->forAll $($ elem $\mid \mathrm{b}$ ag- >count $(\mathrm{elem})=$ bag $2->\operatorname{count}($ elem $))$ and
bag2->forAll(elem | bag 2->count(elem) = bag ->count(elem)) )
bag->union(bag2 : B ag(T)) : B ag(T)
The union of bag and bag. post: result->forAll( elem |
result->count(elem) = bag->count(elem) + bag2->count(elem)) post: bag->f rAll( elem | result->count(elem) = bag->count(elem) + bag2->count(elem)) post: bag2->forAll( elem | result->count(elem) $=$ bag->count(elem) + bag2->count(elem))
bag->union(set : Set (T )) : Bag (T)
The union of bag and set. post: result->forAll(elem | result->count(elem) = bag->count(elem) + set->count(elem)) post: bag->f orAll(elem |
result->count(elem) $=$ bag->count(elem) + set->count(elem)) post: set->forAll(elem |
result->count(elem) $=$ b ag->count(elem) + set $->\operatorname{count}(e l e m))$
bag->intersection(bag2: Bag(T)) : Bag(T)
The intersection of bag and bag2.
post: result->forAll(elem |
result->count(elem) = bag->count(elem).min(bag2->count(elem)) )
post: bag->f orAll(elem |
result->count(elem) = bag->count(elem).min(bag2->count(elem)) )
post: bag2->forAll(elem |
result->count(elem) = bag->count(elem).min(bag2->count(elem)) )
bag->intersection(set : Set(T)) : Set(T)
The intersection of bag and set. post: result->forAll(elem |
result->count(elem) = bag->count(elem).min(set->count(elem)) ) post: bag->f orAll(elem |
result->count(elem) $=$ bag->count(elem).min(set->count(elem)) )
post: set->forAll(elem |
result->count(elem) $=$ bag->count(elem).min(set->count(elem)) )
bag->including(object: T ) : B ag(T)
The bag containing all elements of bag plus object.
post: result->forAll(elem |
if elem $=$ object then
result->count(elem) = bag->count(elem) +1
else
result->count(elem) = bag->count(elem)
endif)
post: bag->f orAll(elem |
if elem = object then
result->count(elem) $=$ bag->count(elem $)+1$
else
result->count(elem) $=$ bag->count(elem $)$
endif)

## bag->excluding(object: T) : Bag (T)

The bag containing all elements of bag apart from all occurrences of object.
post: result->forAll(elem |
if elem $=$ object then
result->count(elem) $=0$
else
result->count(elem) $=$ bag->count(elem $)$
endif)
post: bag->forAll(elem |
if elem $=$ object then
result-> count(elem) $=0$
else
result->count(elem) $=$ bag->count(elem $)$
endif)

## bag->select(expr : OclE expression) : Bag (T )

The sub-bag of bag for which exp is true. post: result $=$ bag->iterate(elem; acc $: \mathrm{Bag}(\mathrm{T})=\mathrm{Bag}\{ \} \mid$
if expr then acc->including(elem) else acc endif)
bag->reject(expr : OclEx pression) : B ag (T)

The sub-bag of bag for which expr is false. post: result $=$ bag->select(not expr)
bag->collect(expr: OclE xpression) : Bag(e xpr.evaluationType)
The Bag of elements which results from applying expr to every member of bag.
post: result $=$ bag->iterate (elem; acc $: B \operatorname{ag}($ expr.evaluationType $)=\mathrm{Bag}\{ \} \mid$ acc->including(expr) )
bag->count(object: T) : Integer
The number of occurrences of object in bag.

## bag->asSequence: Sequence (T)

A Sequence that contains all the elements from bag, in undefined order.
post: result->forAll(elem |bag->count(elem) $=$ result->count(elem)) post: bag->forAll(elem | bag->count(elem) = resu lt->count(elem))

## bag->asSet : Set( T)

The Set containing all the elements from bag, with duplicates re moved.
post: result->forAll(elem | bag->includes(elem) )
post: bag->forAll(elem | result->includes(elem))

### 7.8.2.4 Sequence

A sequence is a collection where the elements are ordered. An element may be part of a sequence more than once.

Properties of Sequence(T), where the instance of Sequence is called sequence.
sequence->count(object : T) : Int ger
The number of occurrences of object in sequence.

```
sequence = (seq uence2 : Sequ ence(T)) : B oolean
```

True if sequence contains the same elements as sequence 2 in the same order.
post: result $=$ Sequence $\{1 .$. sequence- >size $\}->$ forAll(index $:$ In leger $\mid$
sequence->at(index) = sequence2->at(index))
and
sequence->size $=$ sequence $2->$ size

```
sequence->union (sequence2 : Seq uence(T)) : Sequence(T)
```

The sequence consisting of all elements in sequence, followed by all elements in sequence 2 .
post: result- >size $=$ seq uence- $>$ size + seq uence $2->$ size
post: Sequence $\{1$..sequenc e->size $\}->$ forAll(inde $x$ : Ante ger
sequence->at(index) = result->at(index))
post: Sequence $\{1$..sequence2->size $\}->$ forAll(index : Integer $\mid$
sequence2->at(index) =
result->at(index + sequence->size)))
sequence->append (object: T) : Sequence (T)
The sequence of elements, consisting of all elements of sequence, followed by object.
post: result- $>$ size $=$ seq uence- $>$ size +1
post: result->at(result->size) $=$ ob ject
post: Sequ ence $\{1$..sequence-> size $\}->$ forAll(inde $x$ : Ante ger result $->$ at $($ index $)=$ sequence $->$ at( index $)$ )
sequence->prepend(object : T) : Sequ ence(T)
The sequence consisting of object, followed by all elements in sequence.
post: result- >size $=$ seq uence- $>$ size +1
post: result->at(1) = object
post: Sequ ence $\{1$..sequence-> size $\}->$ forAll(inde $x:$ Ante ger $\mid$
sequence->at(index) $=$ result->at(index +1$)$ )

```
sequence->subSequence(lower : In teger, upper : In teger) : Se quence(T)
    The sub-sequence of sequence starting at number lower, up to and including
    element number upper.
    pre : \(1<=\) lower
    pre : lower <= upper
    pre : up per <= seq uence->size
    post: result->size \(=\) u pper -lower +1
    post: Sequence \{lower..upper\}->forAll( in dex |
        result->at(index - lower +1 ) =
        sequence->at(index))
    endif
```

sequence->at(i : Inte ger) : T
The $i$-th element of sequence.
pre : i >=1 and $\mathrm{i}<=$ seq uence->size
sequence->first: T
The first elem ent in sequence.
post: result $=$ sequence->at(1)
sequence->last: T
The last element in sequence.
post: result $=$ sequence->at(sequence->size $)$

## sequence->including(object: T) : Se quence(T)

The sequence containing all elements of sequence plus object added as the last element.
post: result $=$ sequence. $\operatorname{append}($ object $)$

```
sequence->excluding(object : T) : Sequence(T)
```

The sequence containing all elements of sequence apart from all occurrences of object.
The order of the remaining elements is $n$ ot changed.
post:result->includes(object) $=$ false
post: result->size $=$ seq uence->size - sequence->count(object)
post: result = sequence->iterate(elem; acc : Sequ ence(T)
$=$ Sequence $\} \mid$
if elem = object then acc else acc->append(elem) endif )
sequence->select(expression: OclEx pression) : Seq uence(T)
The subsequence of sequence for which expression is true.
post: result $=$ sequence- >iterate $($ elem; acc $:$ Seq uence $(T)=$ Sequence $\} \mid$
if e xpr then acc->including(elem) else acc en if)
sequence->reject(expression : OclEx pression) : Sequ ence(T)
The subsequence of sequence for which expression is false.
post: result $=$ sequence->select(not expr)
sequence->collect(expression : OclExp ression) :
Sequence(expression.evaluationType)
The Sequence of elements which results from applying expression to every member of sequence.
sequence->iterate(expr : OclExp ression) : e xpr.evaluationType
Iterates over the sequence. Iteration will be done from element at position 1 up until the element at the last p osition following the order of the sequ ence.

```
sequence->asBag(): Bag(T)
```

The Bag containing all the elements from sequence, including duplicates.
post: result->forAll(elem | seq uence->count(elem) $=$ result->count(elem) ) post: sequence->forAll(elem $\mid$ seq uence->count(elem) $=$ result- $>\operatorname{count}($ elem $)$ )
sequence->asSet() : Set (T)
The Set containing all the elements from sequence, with duplicated removed.
post: result->forAll(elem | seq uence->includes(elem))
post: sequence->forAll(elem | result->includes(elem))

### 7.9 Grammar

This section describes the grammar for OCL expressions. An executable LL(1) version of this grammar is a vailable on the OCL web site. (See
http://www.software.ibm.com/ad/ocl).
The grammar description uses the EBNF syntax, where "!" means a choice, "?" optionality, and "*" mean s zero or more times, $+m$ means one or more times. In the description of the name, typeName, and string, the syntax for lexical tokens from the JavaCC parser generator is used. (See http://www.suntest.com/JavaCC.)

```
constraint := contextDeclaration
(stereotype name? ":" expression)+
contextDeclaration := "context"
    (classifierContext | operationContext)
classifierContext := (<name> ":") ? <typeName>
operationContext := <typeName> "::" <name>
    "(" formalParameterList? ")"
    ( ":" <typeName> ) ?
formalParameterList := formalParameter (";" formalParameter)*
formalParameter := <name> ":" <typeName>
stereotype := "inv" | "pre" | "post"
expression := letExpression* logicalExpression
ifExpression := "if" expression
    "then" expression
    "else" expression
```

```
    "endif"
logicalExpression := relationalExpression
    ( logicalOperator
        relationalExpression )*
relationalExpression:= additiveExpression
    ( relationalOperator
    additiveExpression ) ?
additiveExpression:= multiplicativeExpression
    ( addOperator
    multiplicativeExpression )*
multiplicativeExpression:= unaryExpression
    ( multiplyOperator unaryExpression )*
unaryExpression := ( unaryOperator postfixExpression )
    | postfixExpression
postfixExpression := primaryExpression ( ("." | "->")
            featureCall )*
primaryExpression := literalCollection
            literal
            | pathName timeExpression? qualifier?
                    featureCallParameters?
            | "(" expression ")"
            | ifExpression
featureCallParameters:= "(" ( declarator )?
                            ( actualParameterList )? ")"
letExpression := "let" <name>
                            ( ":" pathTypeName ) ?
                            "=" expression "in"
literal := <STRING> | <number> | "#" <name>
enumerationType := "enum" "{" "#" <name> ( "," "#" <name>
                            )* "}"
simpleTypeSpecifier:= pathTypeName
                            | enumerationType
literalCollection := collectionKind "{"
    expressionListOrRange? "}"
expressionListOrRange:= expression
    ( ( "," expression )+
    | ( ".." expression )
    )?
```

```
featureCall := pathName timeExpression? qualifiers?
    featureCallParameters?
qualifiers := "[" actualParameterList "]"
declarator := <name> ( "," <name> )*
        ( ":" simpleTypeSpecifier ) ? "|"
pathTypeName := <typeName> ( "::" <typeName> )*
pathName := ( <typeName> | <name> )
        ( "::" ( <typeName> | <name> ) )*
timeExpression := "@" <name>
actualParameterList:= expression ( "," expression )*
logicalOperator := "and" | "or" | "xor" | "implies"
collectionKind := "Set" | "Bag" | "Sequence" |
relationalOperator := "=" | ">" | "<" | ">=" | "<=" | "<>"
addOperator := "+" | "-"
multiplyOperator := "*" | "/"
unaryOperator := "-" | "not"
typeName := ( ["a"_"z"] | ["A"-"Z"] | "_" )
        ( ["a"-"z"] | ["0"-"9"] |
                                    ["A"-"Z"] | "_")*
name := ( ["a"_"z"] | ["A"-"Z"] | "_" )
        (["a"-"z"] | ["0"-"9"] |
            ["A"-"Z"] | "_")*
number := ["0"-"9"] (["0"-"9"])*
string := "'" ( (~["'","\\","\n","\r"])
        | ("\\"
    ( ["n","t", "b", "r","f","\\", "'", "\""]
        | ["0"-"7"] ( ["0"-"7"] ) ?
        | ["0"-"3"] ["0"-"7"] ["0"-"7"]
        )
        )
        ) *
"'"
```

